

How to tackle those difficult problems...

There are two principal approaches to tackling NP-hard problems or other “intractable” problems:

- Q Use a strategy that guarantees solving the problem exactly but doesn’t guarantee to find a solution in polynomial time
- Q Use an approximation algorithm that can find an approximate (sub-optimal) solution in polynomial time

Exact solutions

The exact solution approach includes the strategies:

- Q *exhaustive search* (brute force)
 - useful only for small instances
- Q *backtracking*
 - eliminates some cases from consideration
- Q *branch-and-bound*
 - further cuts down on the search
 - fast solutions for most instances
 - worst case is still exponential

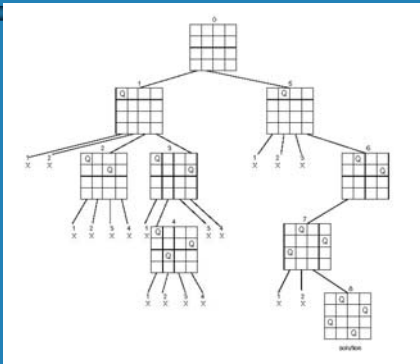
Backtracking

- Q Construct the *state space tree*:
 - nodes: partial solutions
 - edges: choices in completing solutions
- Q Explore the state space tree using depth-first search
- Q “Prune” non-promising nodes
 - dfs stops exploring subtree rooted at nodes leading to no solutions and...
 - “backtracks” to its parent node

Example: The n -Queen problem

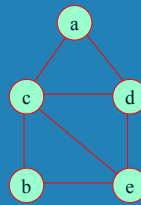
- Q Place n queens on an n by n chess board so that no two of them are on the same row, column, or diagonal

State-space of the four-queens problem



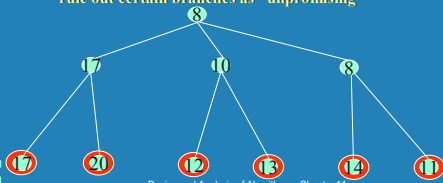
Example: The m -coloring problem

- Q Given a graph and an integer m , color its vertices using no more than m colors so that no two adjacent vertices are the same color.



Branch and Bound

- An enhancement of backtracking
- Applicable to optimization problems
- Uses a lower bound for the value of the objective function for each node (partial solution) so as to:
 - guide the search through state-space
 - rule out certain branches as "unpromising"



Design and Analysis of Algorithms - Chapter 11

7

Example: The assignment problem

- Select one element in each row of the cost matrix C so that:
 - no two selected elements are in the same column; and
 - the sum is minimized

For example:

	Job 1	Job 2	Job 3	Job 4
Person a	9	2	7	8
Person b	6	4	3	7
Person c	5	8	1	8
Person d	7	6	9	4

Lower bound: Any solution to this problem will have total cost of at least:

Design and Analysis of Algorithms - Chapter 11

8

Assignment problem: lower bounds

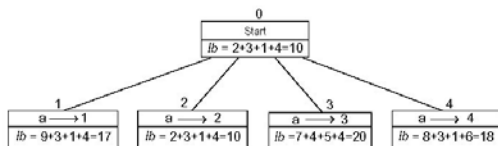


Figure 11.5 Levels 0 and 1 of the state-space tree for the instance of the assignment problem being solved with the best-first branch-and-bound algorithm. The number above a node shows the order in which the node was generated. A node's fields indicate the job number assigned to person a and the lower bound value, lb , for this node.

Design and Analysis of Algorithms - Chapter 11

9

State-space levels 0, 1, 2

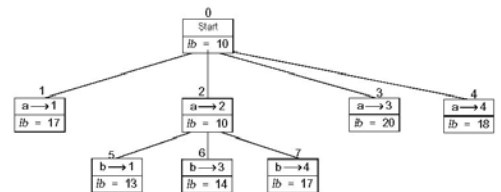


Figure 11.6 Levels 0, 1, and 2 of the state-space tree for the instance of the assignment problem being solved with the best-first branch-and-bound algorithm

Design and Analysis of Algorithms - Chapter 11

10

Complete state-space

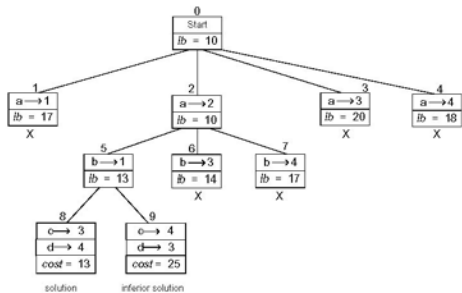
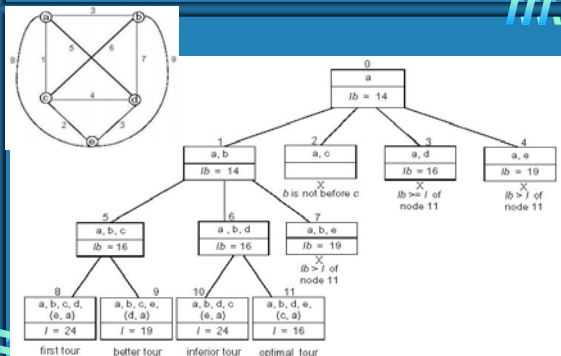


Figure 11.7 Complete state-space tree for the instance of the assignment problem solved with the best-first branch-and-bound algorithm

11

Traveling salesman example:



first tour

better tour

inferior tour

optimal tour